



Time-triggered Message-triggered Object Programming Scheme

K.H. (Kane) Kim
Director of DREAM Laboratory
University of California, Irvine

Juan A. Colmenares
DREAM Lab
University of California, Irvine

*Presented by **Juan A. Colmenares** at the 2006 Monterey Workshop*

October 17th, 2006, Paris, France

Low Productivity Characterizes Real-time Programming

- n **Providing time guarantees** is essential for distributed real-time (DRT) applications
 - n Very complicated if many factors impact response times
- n Conventional practice leads to low productivity
 - n **Avoid most of the software layers**
 - n Implement application software in **C** or **assembly**
- n Increasing demands for challenging new applications that require **precisely timing actions**
 - n Distributed multimedia processing, drive-by-wired, and time-sensitive health care
- n **This practice cannot continue**



Time-triggered Message-triggered Object (TMO) Programming Scheme

- n Initiated in the early 90's
- n A **programmatic approach** for facilitating the development of DRT applications
 - n Contains only **high-level intuitive and yet precise expressions of timing requirements**
 - n From the beginning the objective was to
 - n **Enable design-time guarantees of timely actions**

3



Time-triggered Action

Essence of Real-time Programming

At time T do S

{ = **Start S** during [T - Δ, T + Δ] }

- n **S** may be
 - n A single or compound statement, or a **function** (assume the last one)
- n A **control signal** for activating **S** in a node is derived from the progression of (real) time
 - n A control signal is generated whenever the real-time clock within a node reaches the preset value **T** specified in a scheduling table

4



Time-triggered Objects

Object-oriented Support for Time-triggered Actions

- n Include a new group of member functions
 - n Spontaneous methods (SpMs)
 - n Also called time-triggered methods
 - n SpMs are executed within specified time windows

5



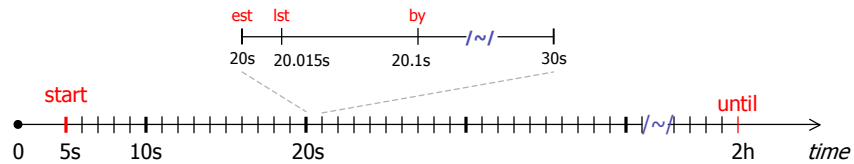
Specification of Execution Time Windows of SpMs

- n Example 1
 - n Start-during (9am, 9:15am)
 - n Finish-by 9:40am
 - n Example 2
 - n For t = from 10am to 10:50am, every 30min
 - n Start-during (t, t+5min)
 - n Finish-by t+10min
 - n It is intuitive and explicit
- } Autonomous Activation Condition (AAC)

6

An AAC Concrete Example

- n MicroSec **from** = 5 * 1000 * 1000; // 5 s
- n MicroSec **until** = 2 * 60 * 60 * 1000 * 1000; // 2 h
- n MicroSec **every** = 1 * 1000 * 1000; // 1 s
- n MicroSec **est** = 0; (early start time)
- n MicroSec **lst** = 15 * 1000; // 15 ms (late start time)
- n MicroSec **by** = 100 * 1000; // 100 ms



Remote Method Calls

Fundamental Mechanism in Distributed Computing

- n Provide transparency in terms of
 - n Location
 - n Low-level communication protocols (e.g., TCP/IP)
- n **Message-triggered object**
 - n Implement remote methods
 - n Also called **service methods (SvMs)**
 - n Activated by messages from clients
- n Real-time application designers must provide **guaranteed completion times** of SvMs

8

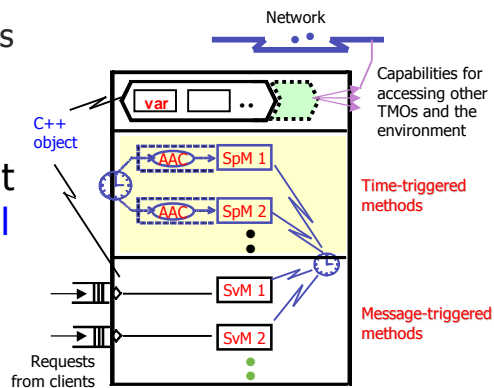
TMO

Time-triggered Message-triggered Object

n Object that also supports

- n Time-triggered actions
- n Remote method calls

n A **natural** and **syntactically small** but **semantically powerful extension** of the conventional object structure



9

First Advantages of the TMO Programming Scheme

n No concerns with

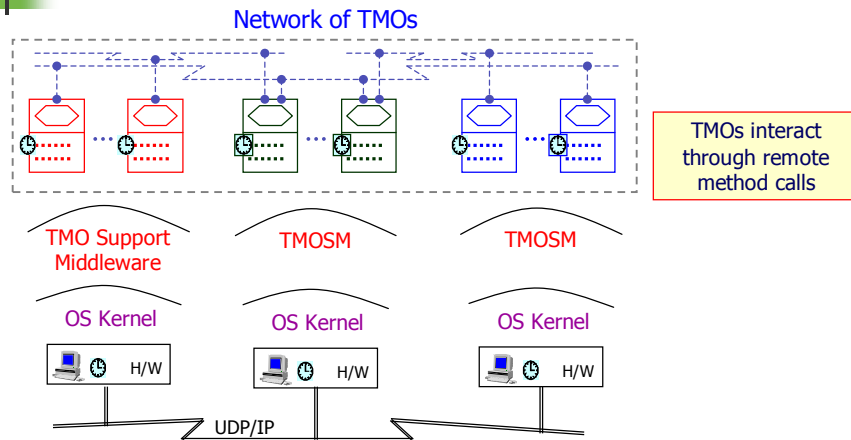
- n Processes and threads
- n Object locations
- n Communication protocols

n No specification of timing requirements in indirect terms

- n e.g., priorities

10

Structure of TMO-based Distributed Real-time (DRT) Application



11

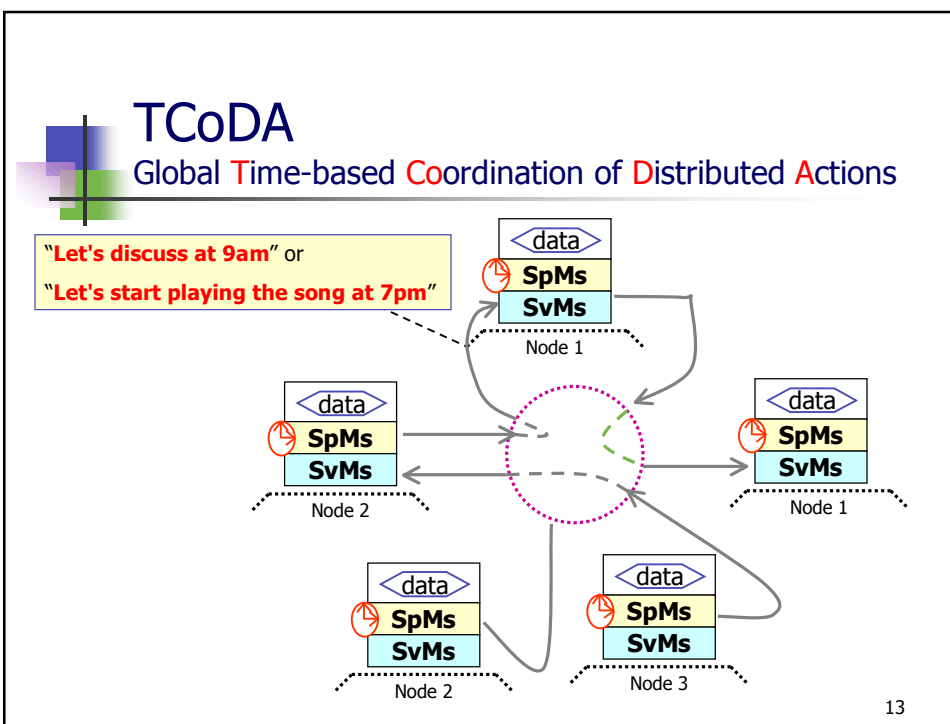
Coordination of Distributed Time-triggered Actions

- n A DRT system should be able to perform **coordinated time-triggered actions** that take place in different nodes

At time T do S ,
 $\{ = \text{Start } S \text{ during } [T - \Delta, T + \Delta] \}$
where $S = [\text{TMO}_0 \text{ does } S_0, \text{TMO}_1 \text{ does } S_1, \dots]$

- n Key requirement
 - n A **global time base** accessible to all nodes
 - n i.e., clock synchronization

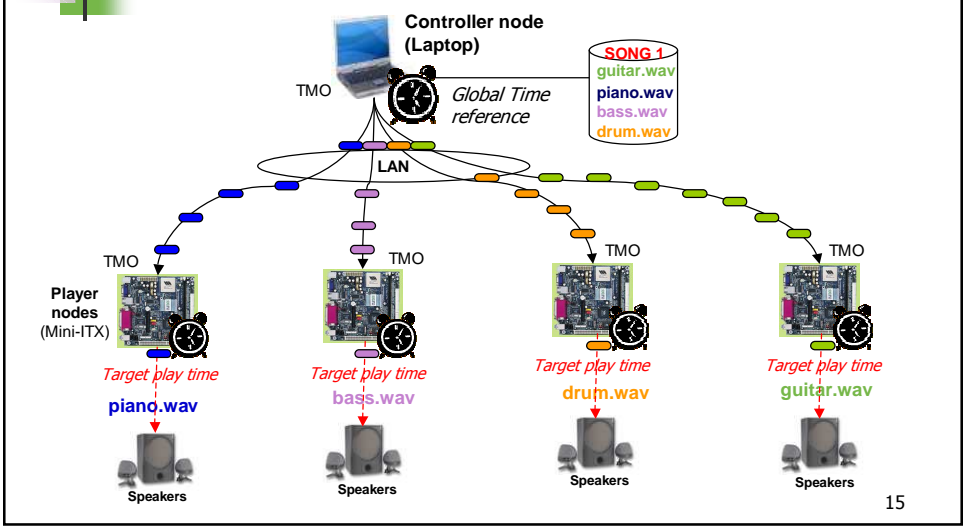
12



- ## TCoDA
- Global Time-based Coordination of Distributed Actions
- n It may **loosen the coupling among subsystems,** and **improve efficiency**
 - n Several nodes can be designed so that they simultaneously start to perform certain actions at 10:00am **without exchanging any message** if they observe certain conditions by 9:59am
 - n That is, less number of messages to exchange between nodes
- 14

Digital Music Ensemble

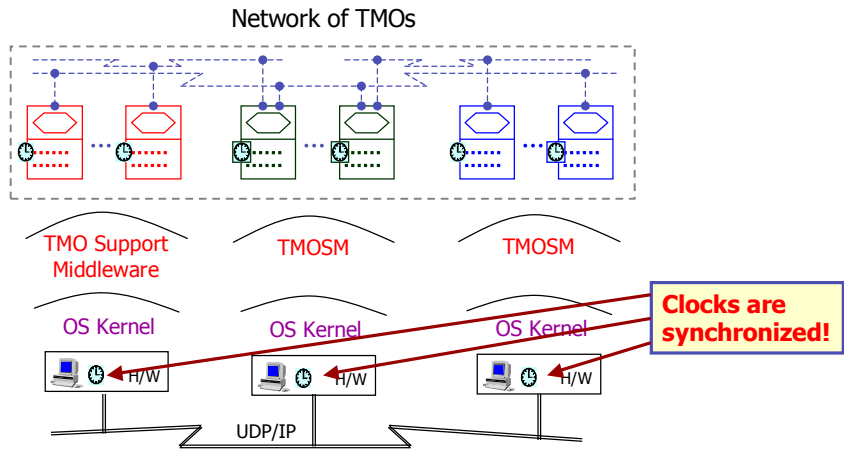
Demonstrating the Power of TCoDA and TMO



TCoDA Has Not Been Sufficiently Explored

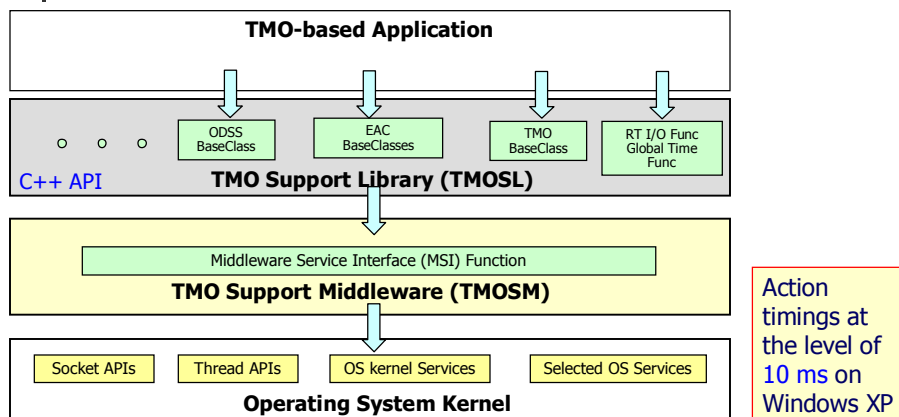
- n **Incorrect perception** about the difficulty of establishing a sufficiently precise global time base
- n **Reality is far better**
 - n 100 μ s-level-precision
 - n Conventional LAN with strict control over the network traffic
 - n 1 μ s-level-precision
 - n **Nodes equipped with GPS receivers** capturing time announcements with microsecond-level accuracy, even if nodes are dispersed over an area larger than a campus

TMO Supports the TCoDA Principle



17

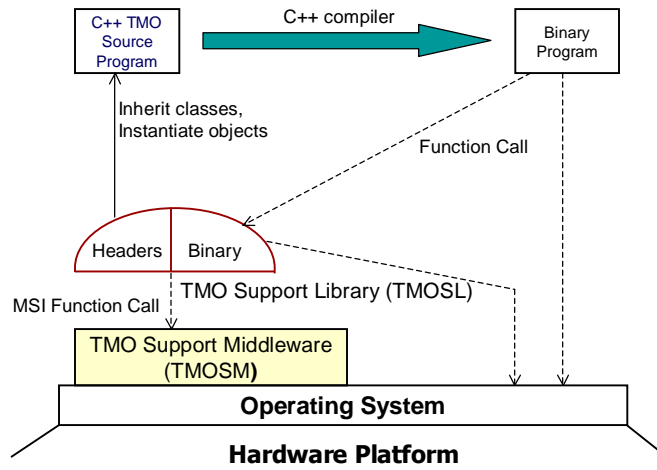
TMOSM and TMOSL



Supported OSs: Windows XP, Windows CE, and Linux 2.6

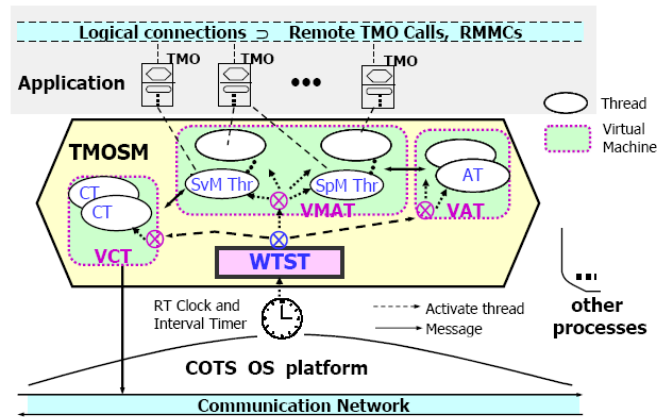
18

TMO Programming Environment



19

TMOSM Architecture



WTST: Watch-dog Timer and Scheduler Thread **VCT:** VM for Communication Threads
VMAT: VM for Main Application Threads **VAT:** VM for Auxiliary Threads

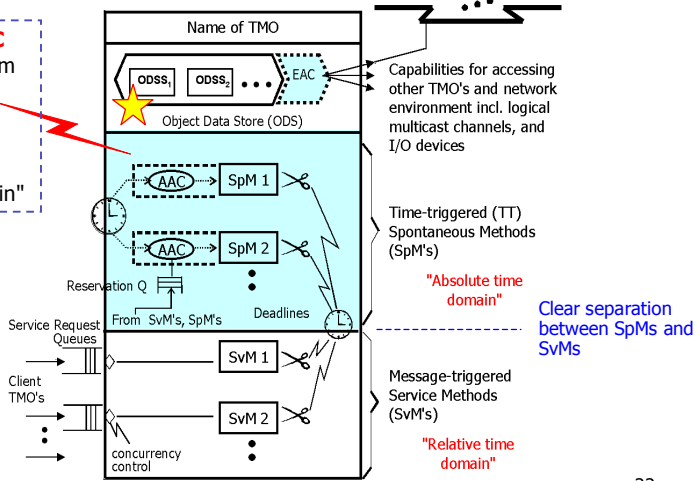
20

A More Detailed View of TMO Programming Scheme

Revisiting Spontaneous Methods

Example of AAC
 "for t = from 10am
 to 10:50am
 every 30min
 start-during
 (t, t+5 min)
 finish-by t+10min"

SpMs must be registered to the execution engine
 Generally done in the constructor





ODSS

Object Data Store Segment

- n A group of data members that represents part of **the internal state** of the enclosing TMO
- n **Basic unit of storage**, which can be
 - n **Locked for exclusive access** by a certain TMO method execution
 - n **Shared** by multiple concurrent executions of TMO methods
- n **Concurrency control**
 - n When a TMO method (SpM or SvM) is invoked, **all the ODSSs** to be accessed during that method execution are **locked** before the execution begins

23



ODSS

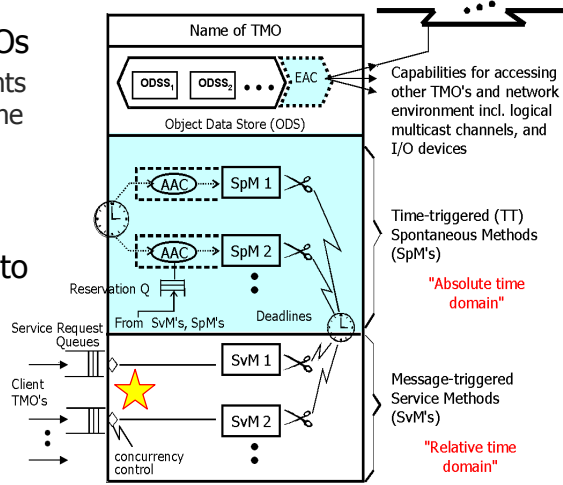
Object Data Store Segment

- n `my_odss.ReleaseODSS ();`
 - n Used when the associated ODSS will not be accessed during the remainder of the method execution
- n Such an early release **enables earlier initiation of other TMO method executions**
- n Once a ODSS is released via `ReleaseODSS ()`, it **cannot be locked again** in the method execution

24

SvMs Service Methods

- n Invoked by client TMOs
 - n Local and remote clients call a SvM exactly in the same way
- n SvMs are ordinary methods of a TMO (sub)class registered to TMOSM



25

Timing Specification of SvMs

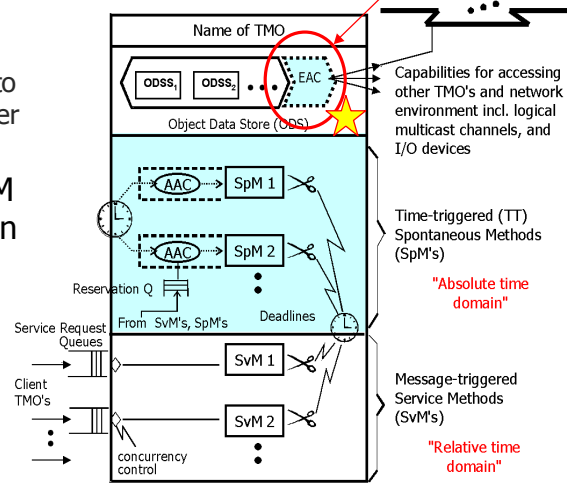
- n The registration of an SvM includes the following parameters:
 - n **Guaranteed execution time bound (GETB)**
 - Optional {
 - n Maximum number of concurrent executions (**PipelineDegree**)
 - n Maximum invocation rate (**MaxInvocations**) and minimum invocation interval (**BasicPeriod**)
 - n Statistical assurances of better service times
- n Other required parameters are:
 - n The names and access modes (RO/RW) of ODSSs
 - n The external name of the SvM
 - n A globally recognized symbolic name

26

Invocation of SvMs

Environment Access Capability: an ODS extension

- n Gate objects
 - n EACs that allow us to invoke SvMs on other TMOs
- n Client's call to a SvM (through a gate) can be associated with **deadline for result return**

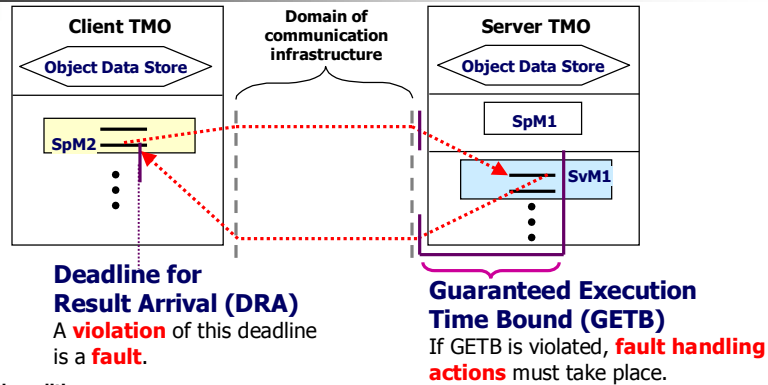


Timing Specification in the Invocations of SvMs

- n `SvMGateClass gate1 (_T("TMO2"), _T("SvM2"), ...);`
- n `gate1.BlockingSR (¶m, sizeof(param), dra, ort);`
 - n `dra = Deadline Result Arrival`
 - n `ort = Official Release Time`
 - n Time at which the invoked SvM should be executed
 - n If 0, SvM is executed ASAP

Toshiba has a patent on the idea of using **ORT** !
 Invented in 1996 and US Patent was granted in April 2001,
 but we learned the invention only in February 2005
 although we have been using it since mid-1999.

Return Deadline vs. Guaranteed Service Time



Required condition

DRA - Call initiation time >
Max. trans. time imposed on comm. infrastructure + **GETB** >
Time consumed by communication infrastructure + **GETB**

29

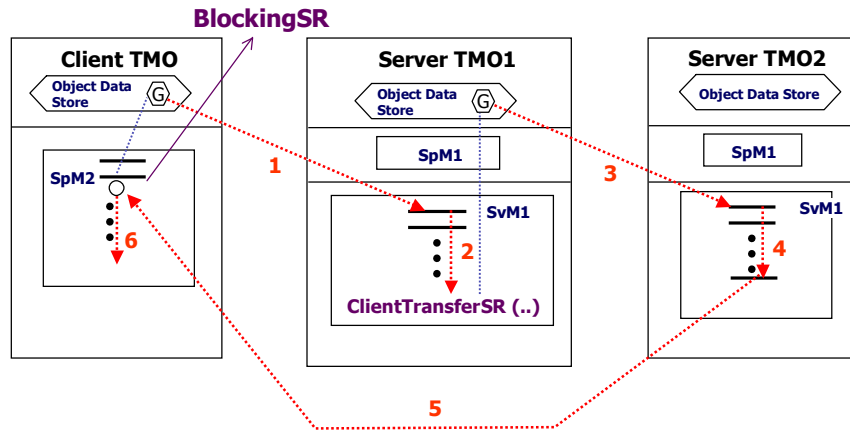
Type of SvM Invocations

- n Blocking calls
 - n with return deadlines imposed and official release times
- n Non-blocking calls and subsequent result checks
 - n with deadlines imposed and official release times
- n One-way calls
 - n With official release times
- n Client-transfer calls
 - n An SvM passes the client's request to another SvM, and the latter returns the result to the client

30



SvM Client-Transfer Call



31



BCC Basic Concurrent Constraint

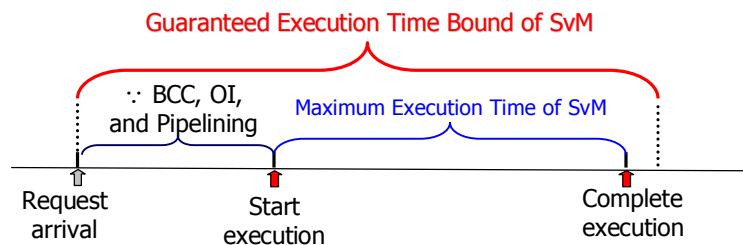
- n SpM executions are given higher priority over SvM executions
 - n An SvM is allowed to execute only if there is no SpM that requires access to the same ODSS and will execute in the time window of this SvM
- n It prevents potential conflicts between SpMs and SvMs and reduces the designer's efforts in guaranteeing timely service capabilities of TMO
 - n Note that this BCC does not impose any restriction on concurrent execution of SpMs or concurrent execution of SvMs
- n Causes potential SvM starvation

32

BCC

Basic Concurrent Constraint

- n In general, the **maximum execution time of an SvM** depends on how many SvMs and non-conflicting SpMs compete for machine resources



33

RMMC

Real-time Multicast and Memory Replication Channel

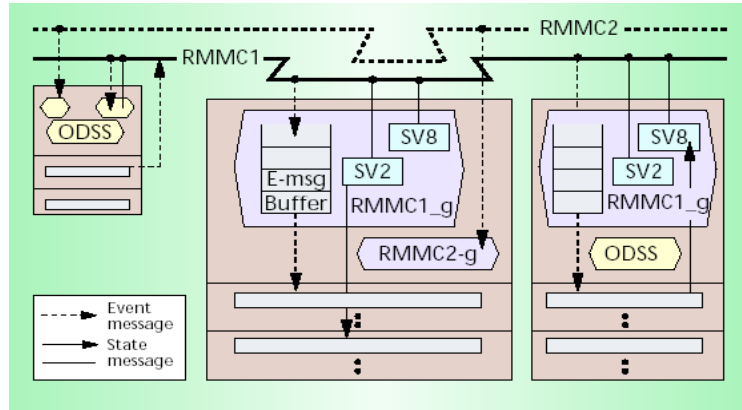
- n TMO also supports
 - n Multicast of event messages
 - n An event message must be read by every corresponding subscriber
 - n Replication of state messages
 - n Consumer are interested in the current state
- n The **producer timestamps the message at message-production time**
- n RMMC also supports **Official Release Time**

34

RMMC

Real-time Multicast and Memory Replication Channel

- n RMMCs accessed by TMO methods
 - n Access gates for 2 RMMCs (RMMC1 and RMMC2) declared in the TMOs



35

Summary

- n TMO programming scheme
 - n Simplifies the development of DRT applications
 - n Allows us to specify the timing requirements of our DRT applications in explicit and intuitive manners
 - n Supports fundamental principles of real-time programming
 - n Time-triggered actions
 - n Global-time-based coordination of distributed actions (TCoDA)

36



Summary

TMO interaction mechanisms	Supported features	Timing specification
Sporadic Methods (SpMs)	Time-triggered actions	AAC { For t = from 10am to 10:50am, every 30min Start-during (t, t+5min), Finish-by t+10min }
Service Methods (SvMs)	Remote methods calls	Server TMO { Guaranteed execution time bound, ...}
		Client TMO { Deadline for Result Arrival, Official Release Time }
Real-time Multicast and Replication Memory Channel (RMMC)	Multicast of event messages and replication of state messages	Producer { Official Release Time }