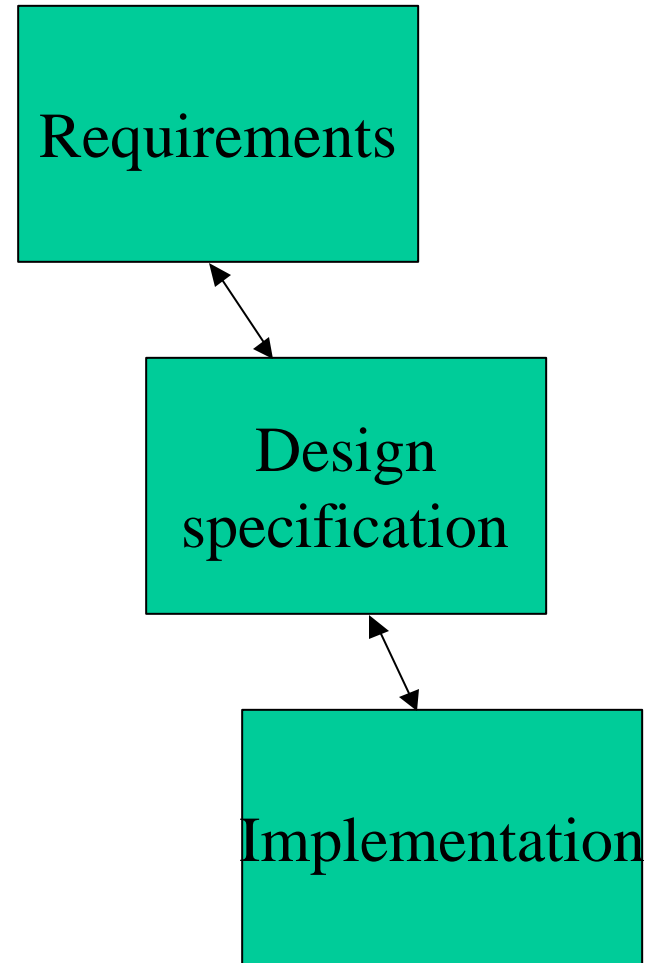

Temporal Logic Based Approach to Test Coverage and Generation

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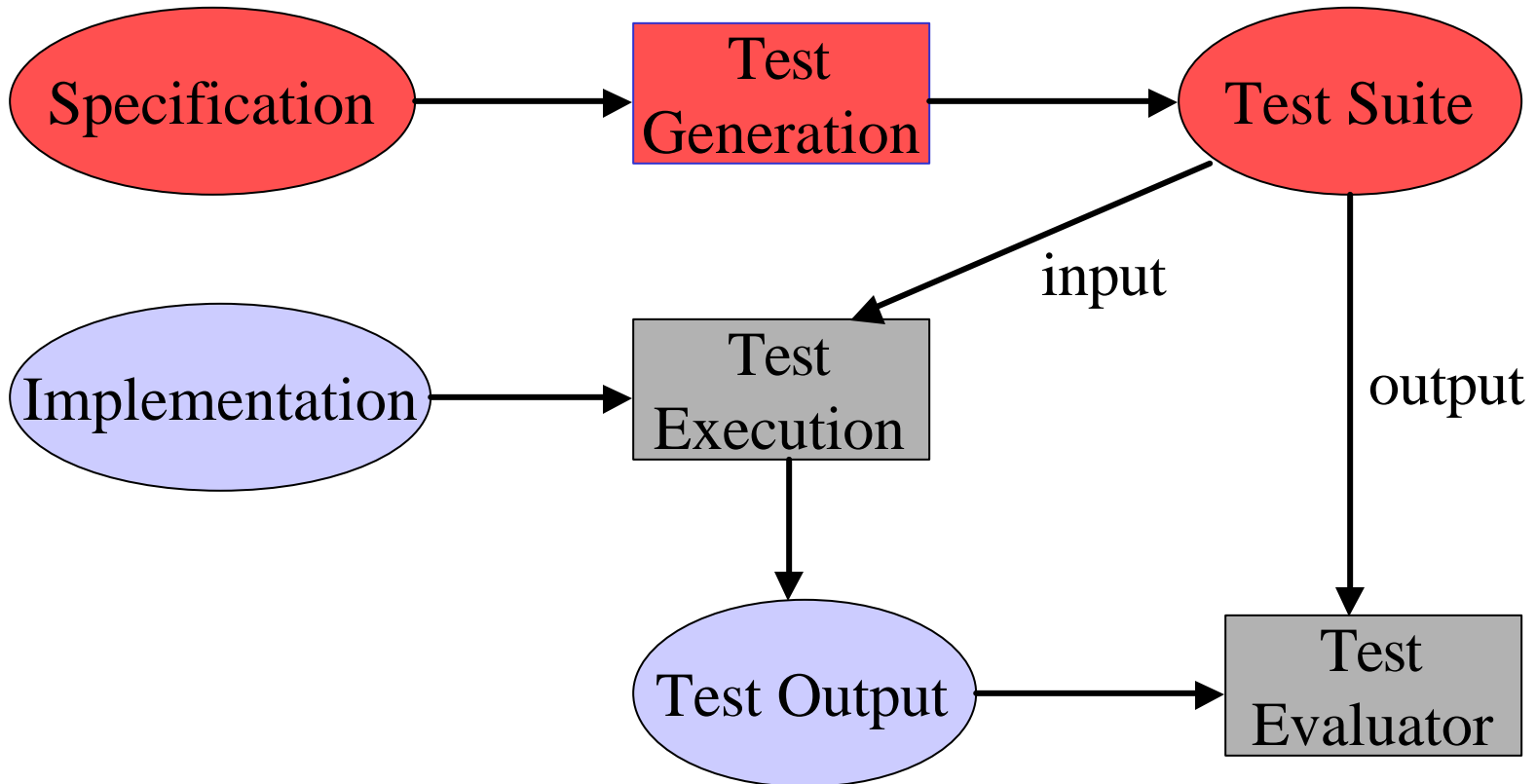
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Software Development Process

- Requirements capture and analysis
 - Informal to formal
 - Consistency and completeness
 - Assumptions and interfaces between system components
 - Application-specific properties
- Design specifications and analysis
 - Formal modeling notations
 - Abstractions
 - Analysis techniques (simulation, model checking, equivalence checking, testing, etc.)
- Implementation
 - Manual/automatic code generation
 - Validation (testing, model extraction, etc.)
 - Run-time monitoring and checking



Overall Structure

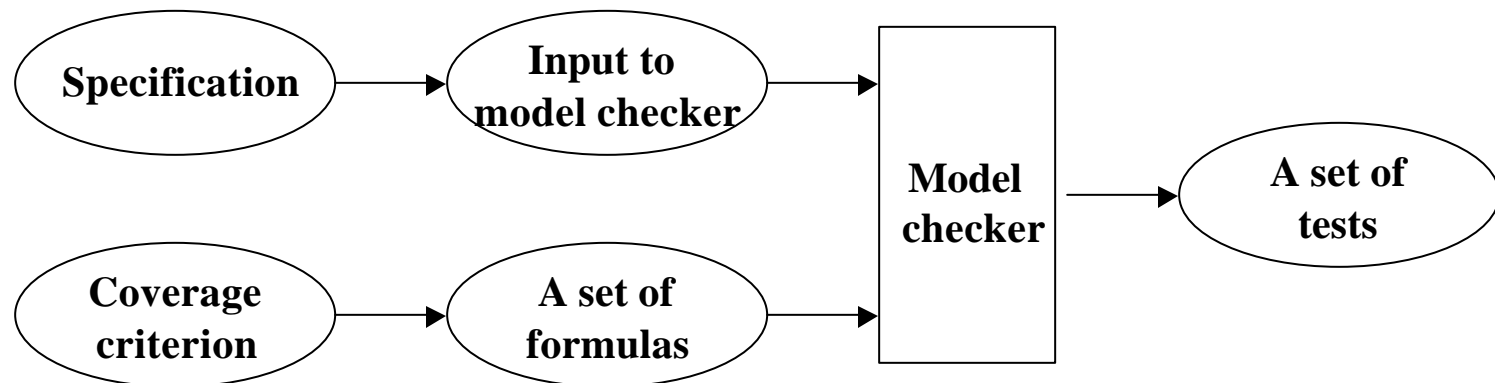


Specification-Based Testing

- Determines whether an implementation conforms to its specification
 - Hardware and protocol conformance testing
 - Widely-used specifications
 - Finite state machines and labeled transition systems
- Two main steps
 - Test generation from specifications
 - What to test, how to generate test
 - Test execution of implementations
 - Applies tests to implementations and validates the observed behaviors

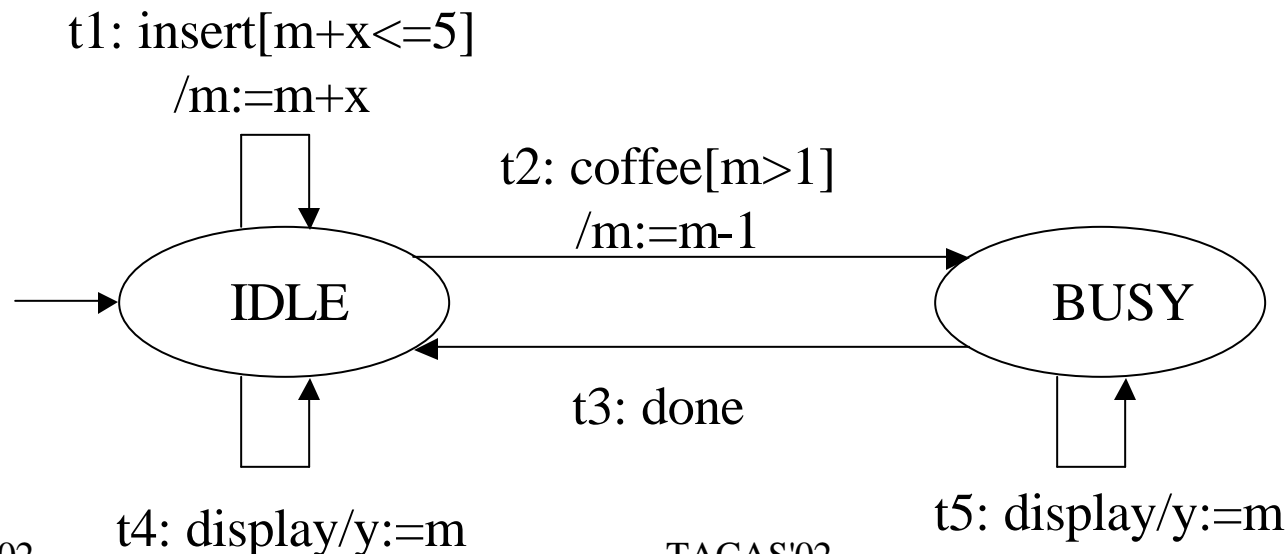
Our Approach

- The problem: automatic test generation from specifications
 - Specifications: EFSMs, Flow graph
 - Coverage criteria
 - Control flow: all-states, all-transitions, etc.
 - Data flow: all-defs, all-uses, all-inputs, all-outputs, etc.
- Test generation



Specifications: EFSM

- $\langle S, S_0, E, V, T \rangle$
 - S : a set of states
 - S_0 : initial state
 - E : a set of input/output events
 - V : variables are manipulated by transitions
 - T : a set of transitions

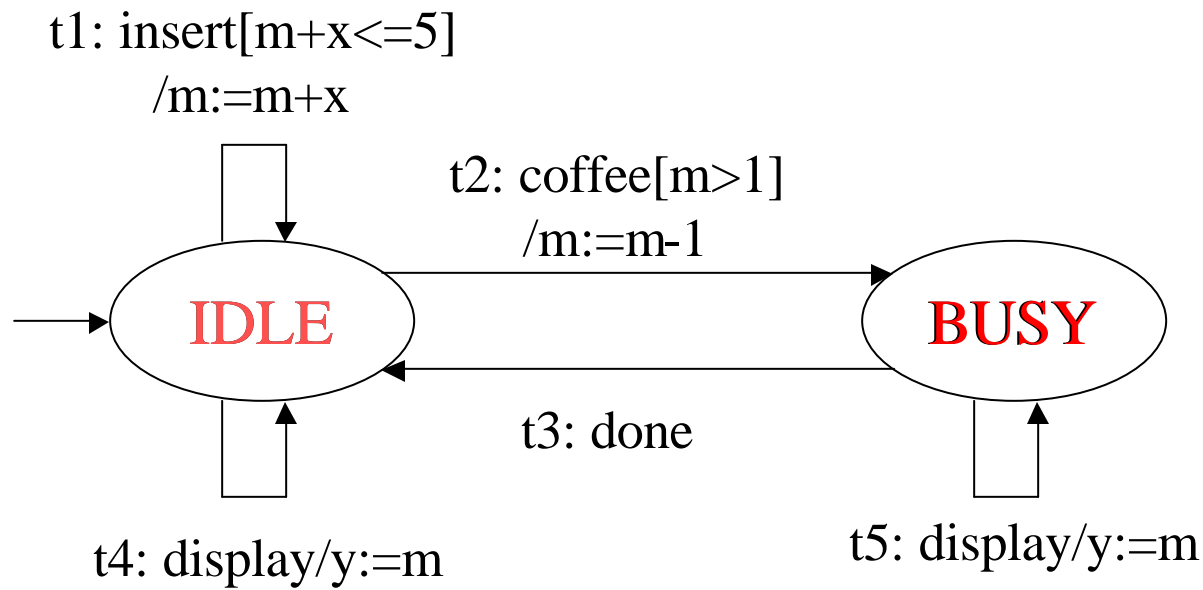


Coverage criteria in WCTL

- Each coverage criterion is represented by a set of temporal logic formulas
 - WCTL: a subset of CTL
 - Atomic propositions p_1, \dots, p_n
 - Temporal operators EX, EU, EF
 - Conjunctions have at most one non-atomic conjuncts
 - Negations can be applied only to atomic propositions
 - Unrestricted disjunctions
 - E.g.: $EF(p_1 \ \& \ EFp_2)$
 - WCTL formulas have linear witnesses

All-states coverage criterion

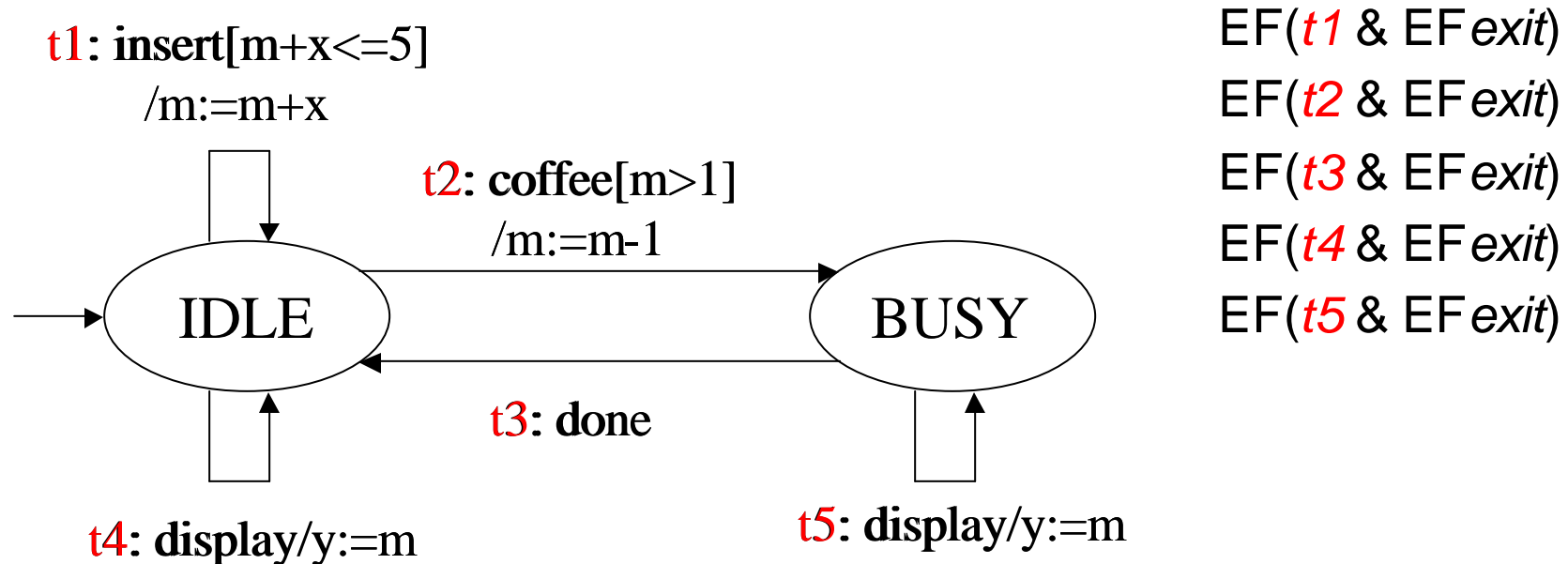
- Requires every state be covered at least once
- With every state s , associate $EF(s)$ & EF_{exit}



$EF(idle)$ & EF_{exit}
 $EF(busy)$ & EF_{exit}

All-transitions coverage criterion

- Requires every transition be covered at least once
- With every transition t , associate $EF(t)$ & EF_{exit}

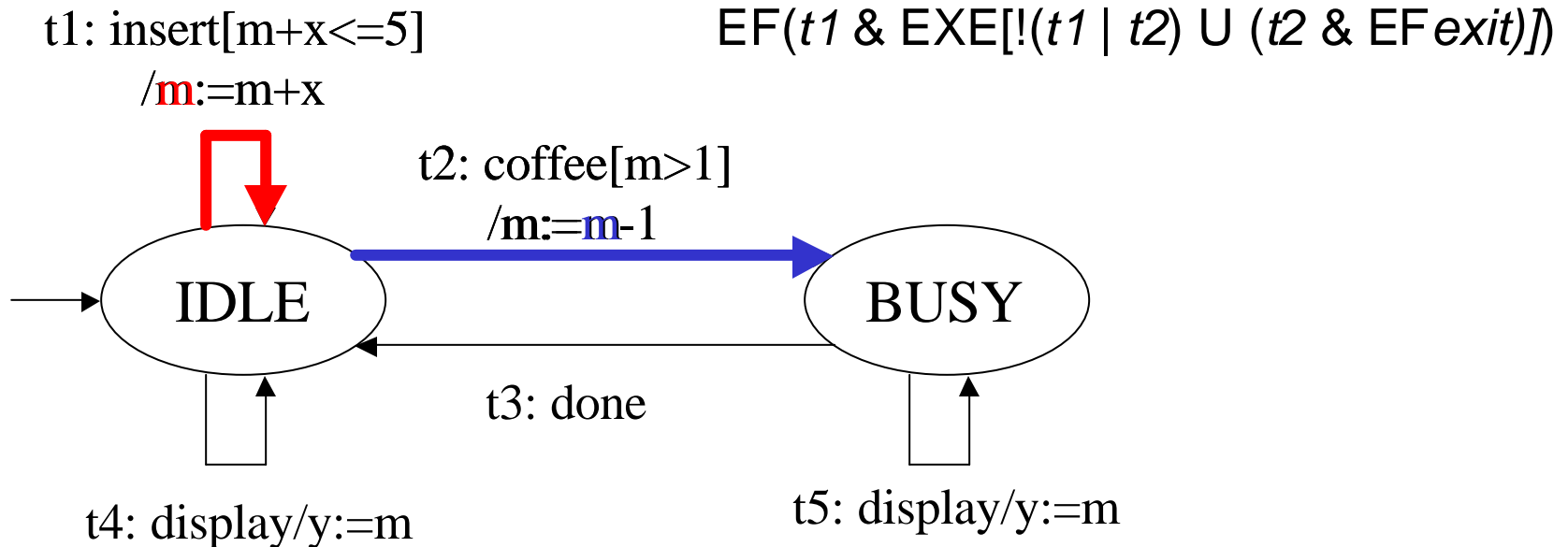


Data flow: definitions and uses

- Central notions in the data-flow analysis
- **Definition**: a value is assigned to a variable
- **Use**: a value of a variable is used in an expression
- Variables are defined and used in transitions
- **Definition-use pair**: (v, t, t')
 - v is defined by t
 - v is used by t'
 - There is a path from t to t' free from other definitions of v

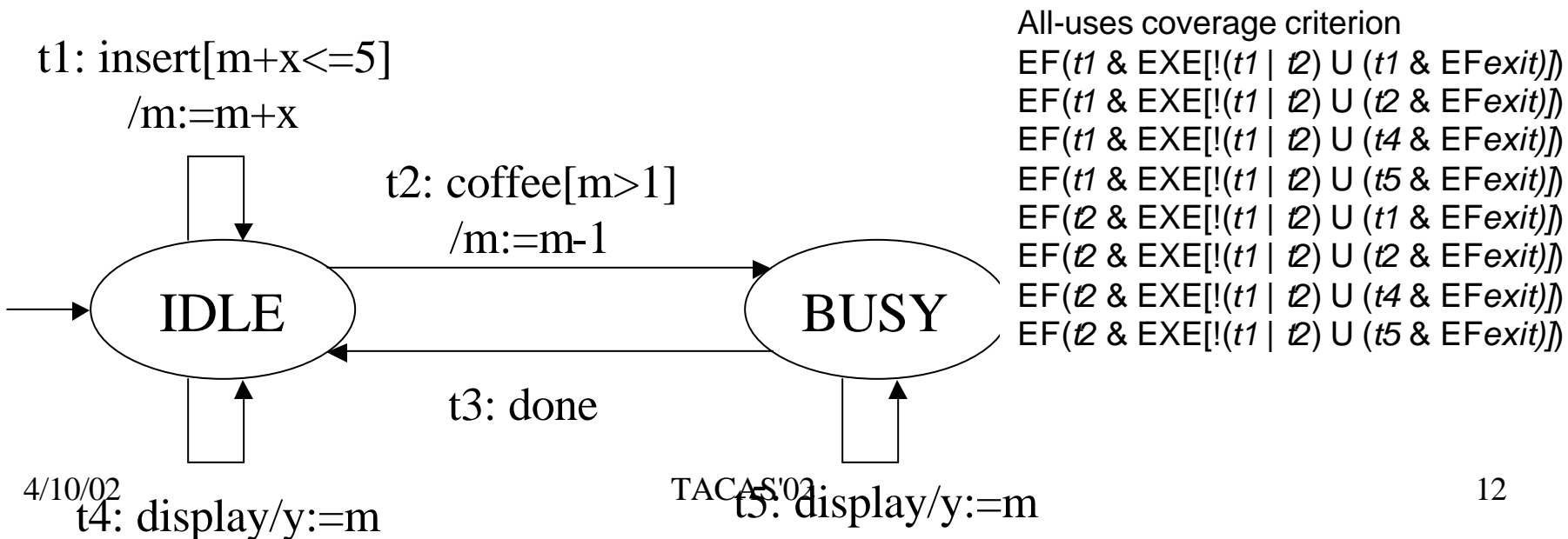
Covering a definition-use pair

- With a definition-use pair (v, t, t') , associate
 - $EF(t \ \& \ EXE[!\text{def}(v) \cup (t' \ \& \ EF_{exit})])$
 - $\text{def}(v)$: disjunction of all transitions that define v



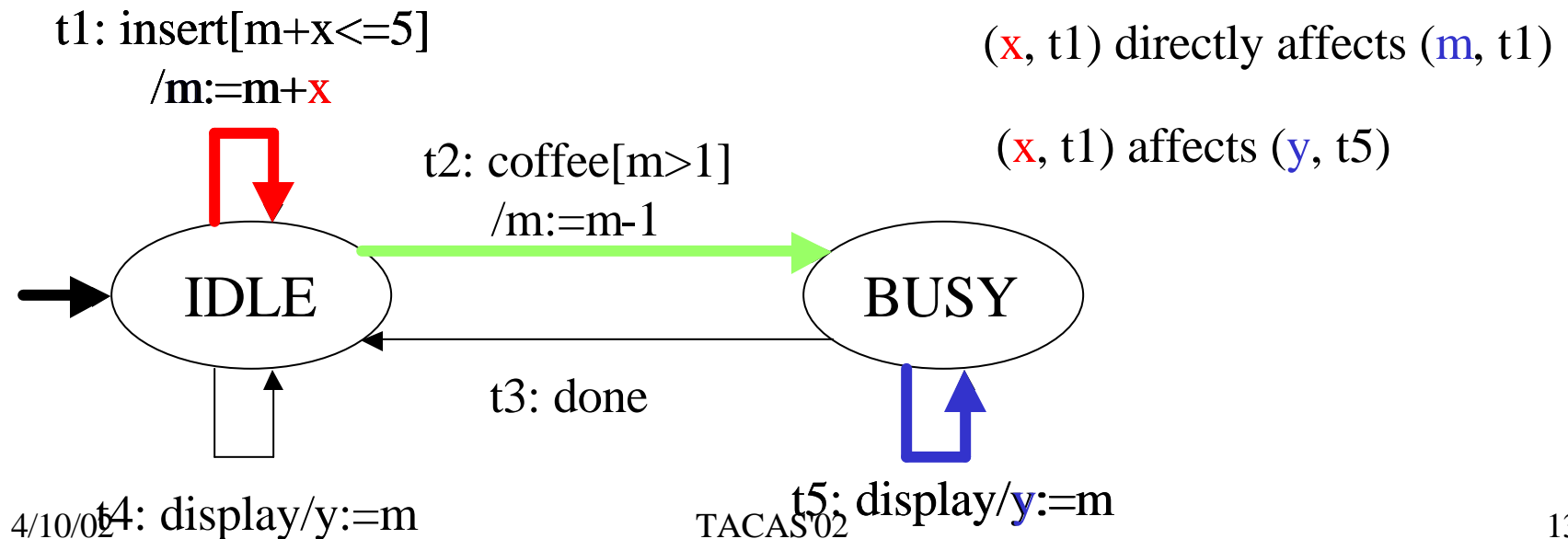
Data-flow coverage criteria

- All-defs coverage criterion
 - Requires a definition-clear path from **every** definition to **some** use be covered at least once
- All-uses coverage criterion
 - Requires a definition-clear path from **every** definition to **every** use be covered at least once



Data flow chains

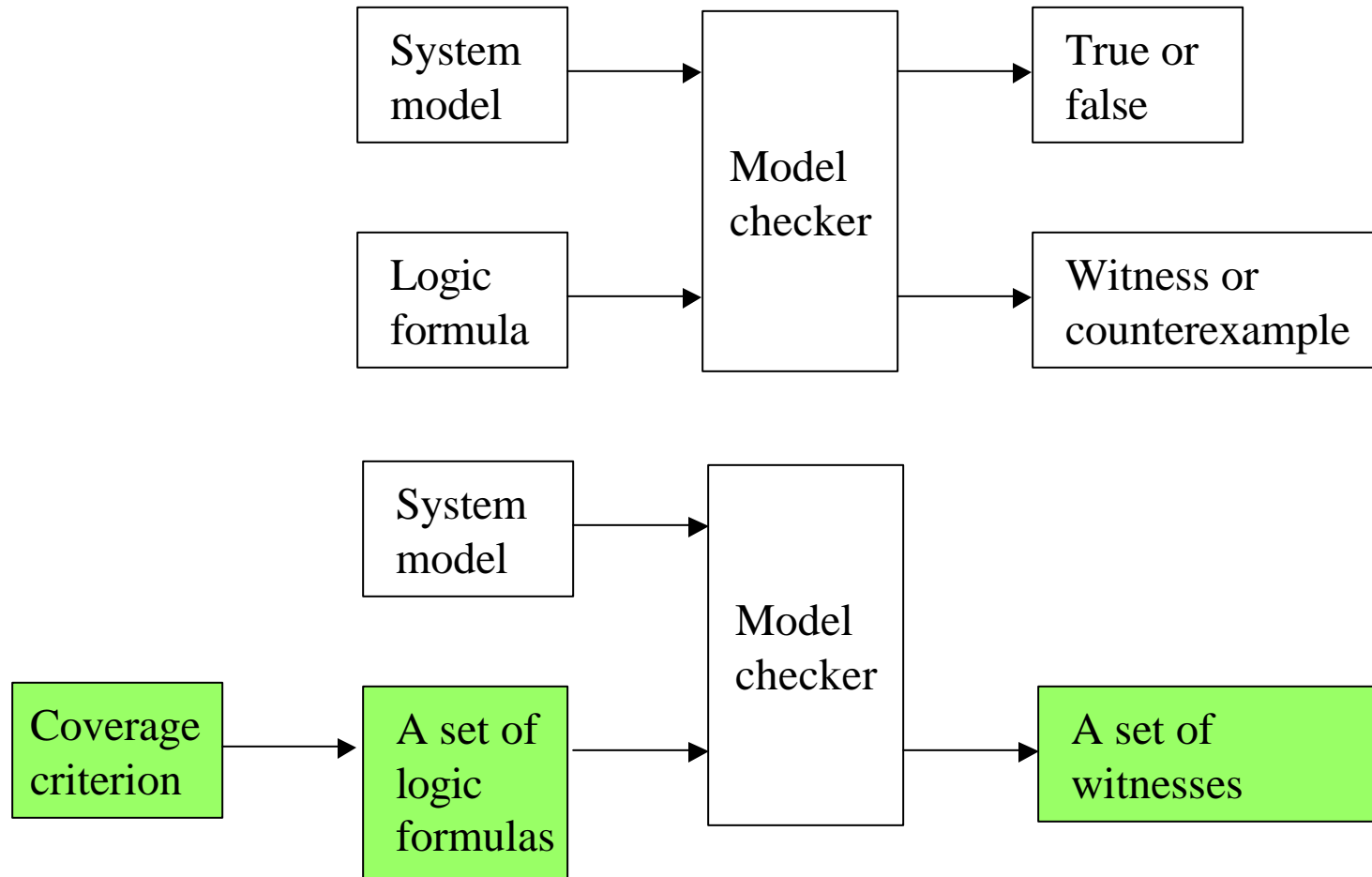
- Affect pair (v, t, v', t')
 - value of v used by t affects the value of v' defined at t'
 - (v, t) directly affects (v', t')
 - Either (v, t) directly affects (v', t') or there is a definition-use pair (v'', t, t') such that (v, t) directly affects (v'', t) and (v'', t) affects (v', t')



Data flow chain coverage

- **Affect pair** (v, t, v', t')
 - May consist of an arbitrary number of definition-use pairs
 - We extend CTL with least fixpoint operators
 - Alternatively, we can use (alternation-free) mu-calculus
- **All-inputs coverage criterion**
 - Requires a path from **every** input to **some** output be covered at least once
- **All-outputs coverage criterion**
 - Requires a path from **every** input to **every** output be covered at least once

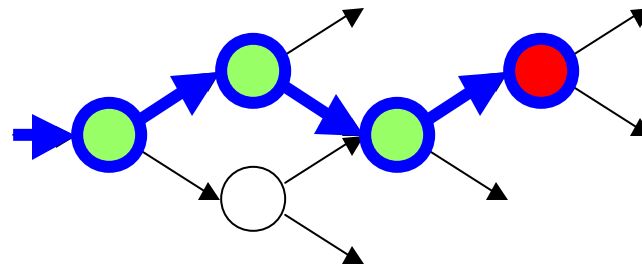
Test Generation



Witness Generation

- Generating a witness for a formula
 - Cost: the length of a witness
 - A minimal-cost witness for a formula
 - Existing model checkers generate a minimal-cost witness by breadth-first search of state space

$E[\text{green} \text{ U } \text{red}]$

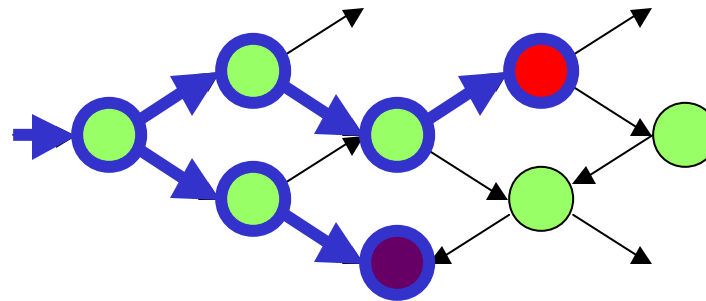


Test Generation

- A set of witnesses for a set of formulas
 - Costs
 - The total length of witnesses or
 - The number of witnesses (reset operation is expensive)
 - Both optimization problems are NP-hard (Hitting Set Problem)

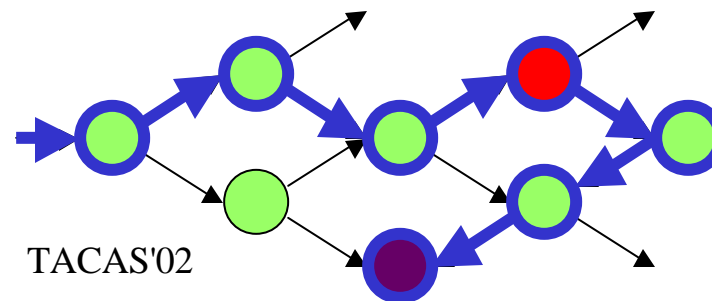
$E[\text{green} \cup \text{red}]$

$E[\text{green} \cup \text{purple}]$



$E[\text{green} \cup \text{red}]$

$E[\text{green} \cup \text{purple}]$

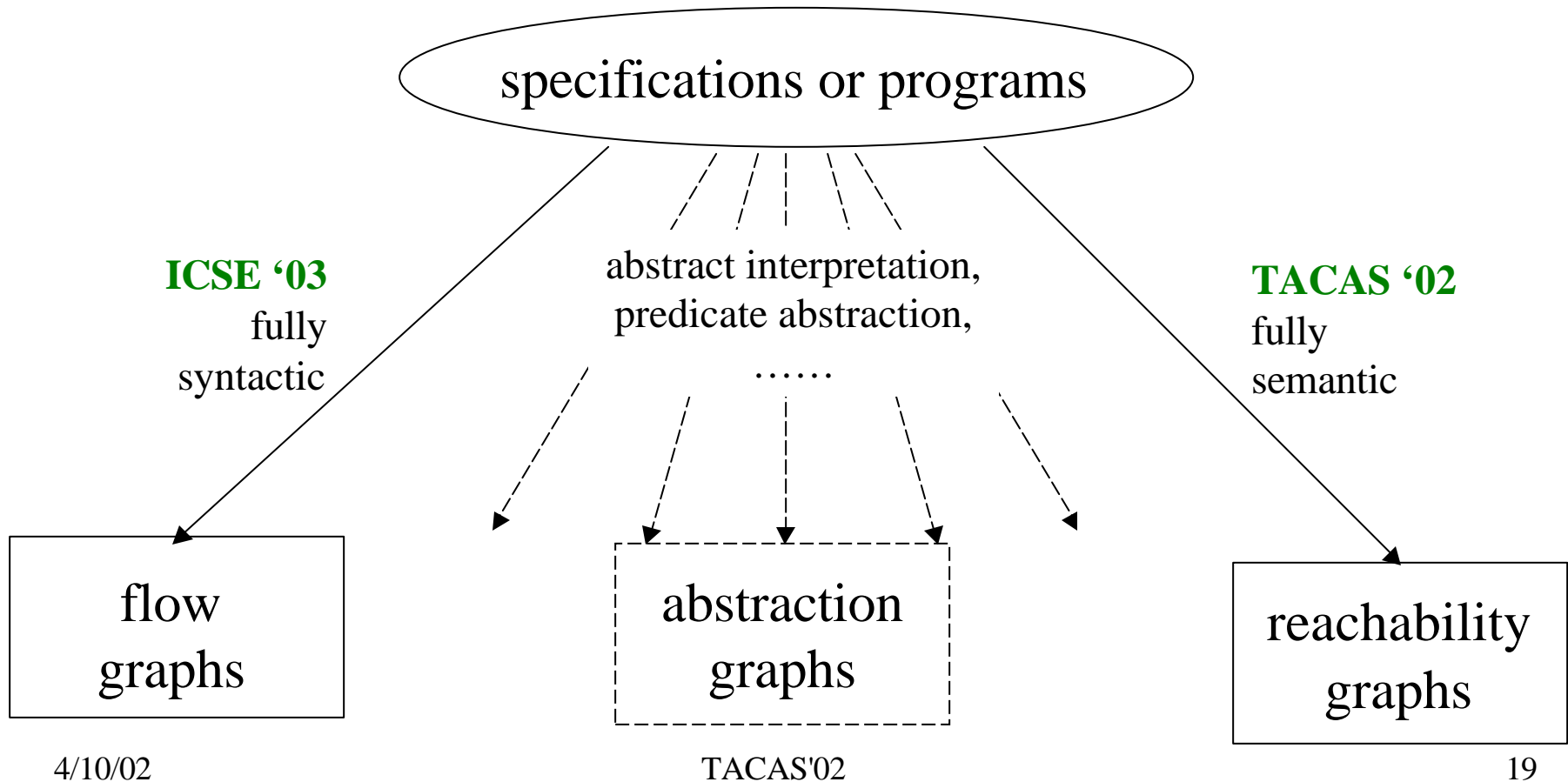


Model-based test generation

- Testing remains the primary validation technique
- Model-based test generation adds rigor to testing:
 - Provide test suits based on a formally verified model
 - Conventional testing coverage criteria applied to the model
- Developed a framework for test generation:
 - Model is Extended Finite-State Machines (EFSM)
 - Offers control-flow (e.g., state coverage, transition coverage) and data-flow (e.g., all-def, all-use coverage) criteria
- Test generation from formal specifications
 - Hierarchical reactive modules
 - EFSM + hierarchy + concurrency
 - Hybrid systems
 - CHARON: EFSM + hierarchy + concurrency + differential equations
- Case studies and tool development
 - Infusion pump system, SONY Aibo, FDA bloodbank policy
 - Application to certification process
- Basis for conformance metrics

Current Work

- A unified framework for test generation
 - Test generation is model checking of abstractions



Q & A
